

① MS115

- We can use truth tables to prove

De Morgan's laws

$$\textcircled{i} \quad \text{not}(P \wedge Q) \equiv \text{not } P \vee \text{not } Q$$

$\nwarrow \text{AND}$ $\swarrow \text{OR}$

$$\textcircled{ii} \quad \text{not}(P \vee Q) \equiv \text{not } P \wedge \text{not } Q$$

$\nwarrow \text{OR}$ $\swarrow \text{AND}$

Let's prove \textcircled{i} :

P	Q	not P	not Q	$P \wedge Q$	$\text{not}(P \wedge Q)$	$\text{not } P \vee \text{not } Q$
T	T	F	F	T	F	F
T	F	F	T	F	T	T
F	T	T	F	F	T	T
F	F	T	T	F	T	T

Hence $\text{not}(P \wedge Q) \equiv \text{not } P \vee \text{not } Q$

↑
logically equivalent to.

\textcircled{ii} as exercise.

- For the tutorial sheet (1), a proposition is logically true if it is true in all cases

e.g. the proposition $P \vee \text{not } P$ is logically true:

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P	$\neg P$	$P \vee \neg P$
T	F	T
F	T	T

Recall: The conditional operator $P \Rightarrow Q$ states that if P is true then Q is also true.

It has truth table

P	Q	$P \Rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

For 2 given propositions P and Q , we will want to prove that $P \Rightarrow Q$ is logically true.

Why? Then we know that if P is true, then Q is true.

How do we show $P \Rightarrow Q$ is logically true for 2 given propositions $P \& Q$?

Strategy: show the case where P is true & Q is false cannot occur.

This gives us 3 methods of argument:

Direct argument: Assume P is true, show that Q is true.

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• Contrapositive argument:

Assume that Q is false, show that P is false

(Recall: the contrapositive of $P \Rightarrow Q$
is the logically equivalent statement
that $\neg Q \Rightarrow \neg P$)

• Proof by contradiction:

Assume P is true and Q is false
and derive a contradiction.

Examples: An integer is a whole
number that is zero, positive
or negative ($0, \pm 1, \pm 2, \dots$)

An integer x is even if there is
an integer k such that

$$x = 2k.$$

An integer y is odd if

$$y = 2n + 1 \text{ for some integer } n.$$

Direct argument to show that
 x and y odd integers $\Rightarrow xy$ is odd

• let $x = 2n+1$ and $y = 2m+1$
for some integers n and m

$$\begin{aligned} \text{Then } xy &= (2n+1)(2m+1) \\ &= 4nm + 2n + 2m + 1 \\ &= 2(2nm + n + m) + 1 \end{aligned}$$

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- Contrapositive argument:
let's show that $x^2 \text{ odd} \Rightarrow x \text{ odd}$
by showing $x \text{ even} \Rightarrow x^2 \text{ even}$

→ let $x = 2k$ for some integer

$$\text{Then } x^2 = (2k)(2k) = 2(2k^2)$$

- Proof by contradiction

→ read a proof of the fact that x is not a fraction if $x^2 = 2$.

or

→ read a proof that there are infinitely many prime numbers.

A statement that contains variables can either be true or false depending on the value of the variables.

Natural application: while ($i < 10$)

These statements are called predicates.

When we fix the value of the variable we have a proposition that is true or false.

e.g. $P(n)$: n is an integer greater than 3
 $P(-1)$ is false, $P(3)$ is false,
 $P(n)$ is true for all $n > 3$

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eg. $P(n)$: there exists an integer n such that $n^2 = 4$

$P(-2)$ true, $P(2)$ true, $P(n)$ false otherwise

eg. $P(n)$: $n^2 \geq 0$ ~~for all integers n~~

→ Does this hold for all integers n ?

Yes. How do we prove this?

More generally, let $P(n)$ be some predicate that is defined for all positive integers n .

How do we show $P(n)$ is true in all cases?

- Induction

Let $P(n)$ be a predicate that is defined for all $n \geq 1$.

Suppose that ① $P(1)$ is true.

and ② for all $n \geq 1$,

$P(n) \Rightarrow P(n+1)$ is true.

Then $P(n)$ is true for all $n \geq 1$.

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Eg. Let's show that

$$P(n) : 1 + 2 + \dots + n = \frac{n(n+1)}{2}$$

is true for all $n \geq 1$.

Induction : ① Let's show $P(1)$ is true:

$$P(1) : 1 = \frac{1(1+1)}{2}$$

$$\text{∴ } 1 = 1$$

$P(1)$ is true.

(this is often called our "base case")

② (Our "inductive step")

To show $P(n) \Rightarrow P(n+1)$
holds for all $n \geq 1$.

Let's take a direct approach:

So, we suppose / assume that $P(n)$ is true

$$\text{i.e. } 1 + 2 + \dots + n = \frac{n(n+1)}{2}$$

Let's show that $P(n+1)$ is also true.

Let's take the L.H.S. of $P(n+1)$:

$$\underbrace{1 + 2 + \dots + n}_{(\text{note: this is the L.H.S. of } P(n))} + n + 1$$

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Thus

$$1 + 2 + \dots + n + n+1$$

$$= \frac{n(n+1)}{2} + n+1 , \text{ as } P(n) \text{ is true}$$

$$= \frac{n(n+1) + 2(n+1)}{2}$$

$$= \frac{(n+1)(n+2)}{2}$$

This is the RHS of $P(n+1)$,
i.e. $P(n+1)$ is true.

Hence $P(n)$ is true for all $n \geq 1$.

Eg. 2 Let's prove that
 $7^n - 1$ is divisible by 6
for all $n \geq 1$.

① Induction
Our $P(n)$ is $6 \mid 7^n - 1$

so $P(1)$ is $6 \mid 7^1 - 1 = 6$: true.

② Assume $P(n)$ is true,
i.e. $6 \mid 7^n - 1$.

let's show $P(n+1)$ is true,

i.e. $6 \mid 7^{n+1} - 1$.

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let's show $7^{n+1} - 1$ is a multiple of 6:

$$7^{n+1} - 1 = 7^n \cdot 7 - 1 \\ = (7^n - 1) \cdot 7 + 6$$

As $P(n)$ is true, $7^n - 1 = 6k$
for some k

Hence

$$7^{n+1} - 1 = (7^n - 1) 7 + 6 \\ = (6k) 7 + 6 \\ = 6(7k) + 6 \\ = 6(7k + 1).$$

Thus $P(n+1)$ is true.

Hence $P(n)$ is true for all $n \geq 1$.

Sets

For us, a set will mean a collection of objects, called elements.

We can list the elements and use curly brackets to show we're dealing with a set,
eg. the set S might be $S = \{1, 2, 4\}$

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We use predicates to describe sets with infinitely many elements:

$$S = \{ 2n-1 \mid n \text{ is a positive integer} \}$$

→ This is the set of odd positive integers which is the set of all integers $2n-1$ such that

n is a positive integer.

i.e. ~~$S = \{ x \mid P(x) \}$~~

is the set of x such that $P(x)$ is true.

Some important sets:

$$\mathbb{N} = \{ 1, 2, 3, \dots \}$$
 is the set of natural numbers

$$\mathbb{Z} = \{ 0, \pm 1, \pm 2, \dots \}$$
 is the set of integers

$$\mathbb{Q} = \left\{ \frac{p}{q} \mid p \text{ is an integer and } q \text{ is a natural number} \right\}$$

is the set of rational numbers

$$\mathbb{R} = \{ \text{all decimal numbers} \}$$

We have $\mathbb{N} \subseteq \mathbb{Z} \subseteq \mathbb{Q} \subseteq \mathbb{R}$,

i.e. \mathbb{N} is contained in \mathbb{Z} , etc.
Also, $\{\}$ or \emptyset is the empty set, the set with no elements

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$$\{ \} = \{ x \mid x \neq x \}$$

We say a set A is a subset of a set B if A is contained in B ,

i.e. $\underline{x \in A} \Rightarrow x \in B$.

If x is an element of A , then x is an element of B

When are two sets equal?

When they have the same elements.

So, we can show two sets A and B are equal by showing $A \subseteq B$ and $B \subseteq A$.

Eg. let $A = \{ n \mid n^2 \text{ is an odd integer} \}$

& $B = \{ n \mid n \text{ is an odd integer} \}$.

Then $A = B$ (Why? look back)

We'll see set operations next.